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Neural Coding and Decoding

Alexander Dimitrov Washington State University

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How does neural activity represent information about stimuli from the environment?

Summary of important features of neural codes







Neural Coding and Decoding

The Problem: Determine a *coding scheme:* How does neural activity represent information about environmental stimuli?

Summary of code feature we consider important:

- Any animal perceives its environment only by observing its own internal representation through neural activity.
- The code must deal with uncertainties introduced by the environment and neural architecture. Activity is by necessity *stochastic* at this finer scale.
- An animal needs to recognize the same object on repeated exposures. Failures at this stage may endanger its well-being. Coding has to be mostly *deterministic* at this level.
- Ecological reasons for uncertainty what if the code **was** deterministic? Somebody else can break it!



Optimal quantization

We define "relevant" as preserving the mutual information I(X;Y) in between stimulus and response.

New Goal: Find the quantizers q^* that minimize

$$D_I = I(X;Y) - I(X_M;Y_N)$$

for fixed M and N. The same as minimizing effective distortion $D_{eff} = -I(X_M, Y_N)$.

Information Theory

The Foundation of the Model

- The structure underlying information theory is a probability measure space (source, random variable). An *expectation E*_X is an integral over the probability measure.
- A signal *x* is produced by a source (r.v.) *X* with a probability *p*(*X*=*x*). A signal *y* is produced by another source *Y* with probability *p*(*Y*=*y*).
- A *communication channel* is a relation between two r.v.'s X and Y. It is described by the conditional probability Q(Y | X).
- Entropy: the uncertainty, or self information of a r. v.

$$H(X) \equiv E_X\left(\log\frac{1}{p(X)}\right) \equiv \int \left(\log\frac{1}{p(X)}\right) p(X) dX$$

Conditional Entropy: the reduced uncertainty of one r.v, if another is observed.

$$H(X | Y) = E_{X,Y}\left(\log\frac{1}{Q(X | Y)}\right)$$

• *Mutual Information:* the amount of information that one r.v. contains about another $I(X;Y) \equiv E_{X,Y}\left(\log \frac{p(X,Y)}{p(X)p(Y)}\right)$

Transmission of information

Consider the encoding process in a probabilistic framework. Information theory makes a few important statements on how messages can be transmitted.

- A signal *x*∈*X* is produced by a source with a probability *p*(*x*). A source is characterized by its entropy *H*(*X*): it can be described completely using no more than *H*(*X*)+1 bits per symbol on the average.
- A channel is a relation between two random variables X and Y. It is completely described by the conditional probability q(y|x). A channel is characterized by its capacity $C = \max I(X;Y)$.
- Sending data through noisy channels: the joint source channel coding theorem. A finite alphabet process with entropy rate H(X) can be transmitted through a channel with capacity C with vanishingly small probability of error iff H(X) < C.







Quantization

A source *Y* can be related to another random variable Y_N through the process of quantization (lossy compression). Y_N is referred to as the reproduction of *Y*. The process is defined by a map $q(Y_N|Y): Y \rightarrow Y_N$

called a *quantizer*. In general, quantizers can be stochastic: *q* assigns to $y \in Y$ the probability that the response y belongs to an abstract class $y_N \in Y_N$. A deterministic quantizer (simple function) is a special case in which *q* takes the values of 0 or 1 only. It can be shown that the mutual information I(X; Y) is the least upper bound of $I(X_M; Y_N)$ over all possible reproductions $(X_M; Y_N)$ of (X; Y). Hence, the original mutual information can be approximated with arbitrary precision using carefully chosen reproduction spaces.

Rate Distortion Theory

Rate distortion theory is concerned with reduced representations of random variables (lossy compression). The quality of reproduction (fidelity) is assessed through a *distortion function*.

Consider the quantization $X \rightarrow X_N$.

Definition 3. *A* (pointwise) distortion function, or distortion measure is a mapping

$d: X \times X_N \to R^+$

from the set of source/reproduction pairs into the set of nonnegative reals. The distortion is a measure of the "error" made by representing the symbol $x \in X$ with $x_n \in X_N$. **Example 3** (Squared error distortion). $d(x; x_n) = (x - x_n)^2$. **Definition 4 (Expected (mean) distortion function).** $D(X;X_N) = E_{p(x;xn)}d(x;x_n)$

Definition 5 (Rate distortion problem). The information ratedistortion function R(D) for a source X with a distortion measure $d(x; x_n)$ is defined as

$$R(D) = \min_{q(x_n|x):D(X;X_N) \le D} I(X;X_N)$$

where the minimization is over all conditional probabilities $q(x_n|x)$ for which the joint distribution $p(x; x_n) = q(x_n|x)p(x)$ satisfies the expected distortion constraint. Equivalently, one may consider the distortion-rate problem

$$D(R) = \min_{q(x_n|x):I(X;X_N) \leq R} D(X;X_N)$$

Quantization theory.

The quantized information quantities in Y_N are (Gray '94)

$$I(X_{M};Y_{N}) \equiv \sum_{i,j}^{M,N} p(x_{i},y_{j}) \ln \frac{p(x_{i},y_{j})}{p(x_{i})p(y_{j})} \quad H(X_{M}) \equiv \sum_{i=1}^{M} p(x_{i}) \ln \frac{1}{p(x_{i})}$$

If a quantizer h refines f then

$$Y \xrightarrow{h} Y_{M} \xrightarrow{J} Y_{N} \quad (\text{Markov relation})$$
$$H(Y_{M}) \geq H(Y_{N})$$
$$H(X_{L} | Y_{M}) \leq H(X_{L} | Y_{N})$$
$$I(X_{L}; Y_{M}) \geq I(X_{L}; Y_{N})$$

 $I(X;Y) = \sup I(X_M;Y_N).$

The way we build up measures in measure spaces.





The sensory system challenge: dealing with complex stimuli.

- $I(X_M, Y_N)$ cannot be estimated directly for high-dimensional stimulus sets P(X, Y) not known. Use a model. This produces an upper bound to the distortion (Data processing theorem). **Better model = tighter bound.**
- Here we use a Gaussian estimate of the stimulus:

$$\mathcal{P}_{est}(x \mid m) = N(x; \mu_m, C_m)$$

where (μ_m, C_m) are the stimulus mean and covariance, conditioned on class *m*. The estimate of the stimulus probability then is the Gaussian mixture model

$$p_{est}(x) = \sum_{m=1..M} p_m N(x; \mu_m, C_m)$$

The parametric quantizer is $p(m \mid x) = \frac{p_m N(x; m_m, C_m)}{p_{est}(x)}$.

• *!!!* This model imposes a distance on the input space: it defines when stimuli are close to each other.



















Discussion

- model a set of neurons as a communication channel.
- define a coding scheme through equivalence classes of stimulus/response pairs.
 - Coding is probabilistic on codewords.
 - Coding is almost deterministic on codeword classes.
 - The number of classes is ~ $2^{I(X,Y)}$.
- propose a new method to quantify neural spike trains.
 - Quantize the response patterns to a smaller space.
 - Use an information-based distortion measure.
 - Minimize the information distortion for a fixed size reproduction.
- present results with cricket sensory data.
 - Use temporal patterns of spikes across a few neurons.
 - Recover the stimulus reconstruction kernel at the coarsest quantization.
 - Demonstrate the presence of additional structure at finer quantizations.
 - Demonstrate non-linear processing in several cells.

