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#### Neural Coding and Decoding

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How does neural activity represent information about stimuli from the environment?

# Summary of important features of neural codes





Given a stimulus, we observe many slightly different neural responses (*noisy spike trains*):





## Neural Coding and Decoding

**The Problem:** Determine a *coding scheme:* How does neural activity represent information about environmental stimuli?

#### **Summary of code feature we consider important:**

- Any animal perceives its environment only by observing its own internal representation through neural activity.
- The code must deal with uncertainties introduced by the environment and neural architecture. Activity is by necessity *stochastic* at this finer scale.
- An animal needs to recognize the same object on repeated exposures. Failures at this stage may endanger its well-being. Coding has to be mostly *deterministic* at this level.
- Ecological reasons for uncertainty what if the code **was** deterministic? Somebody else can break it!



## Optimal quantization

We define "relevant" as preserving the mutual information *I(X;Y)* in between stimulus and response*.*

**New Goal:** Find the quantizers *q\** that minimize

$$
D_i = I(X; Y) - I(X_M; Y_N)
$$

for fixed M and N. The same as minimizing effective distortion  $D_{\text{eff}}$  =  $-(X_M, Y_N)$ .

### **Information Theory**

**The Foundation of the Model**

- The structure underlying information theory is a probability measure space (source, random variable). An *expectation*  $E<sub>x</sub>$  is an integral over the probability measure. probability
- A signal *x* is produced by a source (r.v.) *X* with a probability *p(X=x)*. A signal *y* is produced by another source *Y* with probability *p(Y=y).*
- A *communication channel* is a relation between two r.v.'s *X* and *Y*. It is described by the conditional probability *Q(Y | X).*
- **Entropy:** the uncertainty, or self information of a r. v.

$$
H(X) \equiv E_X \left( \log \frac{1}{p(X)} \right) \equiv \int \left( \log \frac{1}{p(X)} \right) p(X) dX
$$

**Conditional Entropy:** the reduced uncertainty of one r.v, if another is observed. ⎞

$$
H(X|Y) \equiv E_{X,Y} \bigg( \log \frac{1}{Q(X|Y)} \bigg)
$$

⎝

• *Mutual Information:* the amount of information that one r.v. contains about another  $\sqrt{2}$ ⎞  $\overline{\phantom{a}}$  $\big($ ≡  $I(X;Y) \equiv E_{X,Y} \bigg( \log \frac{p(X,Y)}{p(X)p(Y)} \bigg)$ 

⎠

#### Transmission of information

Consider the encoding process in a probabilistic framework. Information theory makes a few important statements on how messages can be transmitted messages can be transmitted.

- A signal *x*∈*X* is produced by a source with a probability *p(x)*. A source is characterized by its entropy *H(X):* it can be described completely using no more than *H(X)*+1 bits per symbol on the average.
- A channel is a relation between two random variables *X* and *Y*. It is completely described by the conditional probability q(y|x). A channel is characterized by its capacity  $C = \max_{p(x)} I(X;Y)$ .
- Sending data through noisy channels: the joint *p x*  $\bullet$  Sending data through noisy channels: the joint source channel coding theorem. *A finite alphabet process with entropy rate* H(X) *can be transmitted through a channel with capacity* C *with vanishingly small probability of error iff*  $H(X) < C$ .







## **Quantization**

A source *Y* can be related to another random variable  $Y_N$  through the process of *quantization (lossy compression).*  $Y_N$  *is referred* to as the *reproduction of Y . The process is defined by a map*  $q(Y_N|Y) : Y \to Y_N$ 

called a *quantizer*. In general, quantizers can be stochastic*: q* assigns to  $y \in Y$  *the probability* that the response y belongs to an abstract class  $y_N \in Y_N$ . A deterministic quantizer (simple function) is a special case in which *q* takes the values of 0 or 1 only. It can be shown that the mutual information  $I(X; Y)$  is the least upper bound of  $I(X_M; Y_N)$  over all possible reproductions  $(X_M; Y_N)$  of  $(X; Y)$ . Hence, the original mutual information can be approximated with arbitrary precision using carefully chosen reproduction spaces.

## Rate Distortion Theory

Rate distortion theory is concerned with reduced representations of random variables (lossy compression). The quality of reproduction (fidelity) is assessed through a *distortion function.*

Consider the quantization  $X \to X_N$ .

**Definition 3**. *A (pointwise) distortion function, or distortion measure is a mapping*

$$
d: X \times X_N \to R^+
$$

*from the set of source/reproduction pairs into the set of nonnegative reals. The distortion is a measure of the "error" made by representing the symbol*  $x \in X$  *with*  $x_n \in X_N$ . **Example 3** (Squared error distortion).  $d(x; x_n) = (x - x_n)^2$ .

**Definition 4 (Expected (mean) distortion function).**  $D(X;X_N) = E_{p(x;xn)}d(x;x_n)$ 

**Definition 5 (Rate distortion problem).** The information rate*distortion function R(D) for a source X with a distortion measure*  $d(x; x_n)$  is defined as

$$
R(D) = \min_{q(x_n|x):D(X;X_N)\le D} I(X;X_N)
$$

*where the minimization is over all conditional probabilities*   $q(x_n|x)$  for which the joint distribution  $p(x; x_n) = q(x_n|x)p(x)$ *satisfies the expected distortion constraint. Equivalently, one may*  consider the distortion-rate problem

$$
D(R) = \min_{q(x_n|x): I(X;X_N) \le R} D(X;X_N)
$$

### Quantization theory.

The quantized information quantities in  $Y_N$  are (*Gray '94*)

$$
I(X_M; Y_N) = \sum_{i,j}^{M,N} p(x_i, y_j) \ln \frac{p(x_i, y_j)}{p(x_i) p(y_j)} \quad H(X_M) = \sum_{i=1}^{M} p(x_i) \ln \frac{1}{p(x_i)}
$$

If a quantizer *h* refines *f* then

$$
Y \xrightarrow{h} Y_M \xrightarrow{f} Y_N \quad \text{(Markov relation)}
$$
\n
$$
H(Y_M) \geq H(Y_N)
$$
\n
$$
H(X_L | Y_M) \leq H(X_L | Y_N)
$$
\n
$$
I(X_L; Y_M) \geq I(X_L; Y_N)
$$

 $I(X;Y) = \sup I(X_M;Y_N)$ .

The way we build up measures in measure spaces.





### The sensory system challenge: dealing with complex stimuli.

- $I(X_M; Y_N)$  cannot be estimated directly for high-dimensional stimulus sets *P(X,Y)* not known. Use a model. This produces an upper bound to the distortion (Data processing theorem). **Better model = tighter bound.**
- Here we use a Gaussian estimate of the stimulus:

$$
\mathcal{D}_{\text{est}}(x \mid m) = N(x; \mu_m, C_m)
$$

 $p_{_{est}}(x | m) = N(x; \mu_m, C_m)$  where  $(\mu_m, C_m)$  are the stimulus mean and covariance, conditioned on class *m.* The estimate of the stimulus probability then is the Gaussian mixture model

$$
p_{est}(x) = \sum_{m=1..M} p_m N(x; \mu_m, C_m)
$$

The parametric quantizer is  $p(m | x) = \frac{p_m N(x; m_m, C_m)}{C}$  $(x) = \frac{p_m N(x; m_m, C_m)}{p_{est}(x)}$ .  $p(m | x) = \frac{P_m(x | x, m_m)}{\sqrt{x}}$ *est*  $=\frac{P_m N\left(\lambda, m_m, \mathbf{C}_m\right)}{2}$ 

• *!!!* **This model imposes a distance on the input space: it defines when stimuli are close to each other.**



















#### **Discussion**

- model a set of neurons as a communication channel.
- define a coding scheme through equivalence classes of stimulus/response pairs.
	- Coding is probabilistic on codewords.
	- Coding is almost deterministic on codeword classes.
	- The number of classes is  $\sim 2^{I(X,Y)}$ .
- propose a new method to quantify neural spike trains.
	- Quantize the response patterns to a smaller space.
	- Use an information-based distortion measure.
	- Minimize the information distortion for a fixed size reproduction.
- present results with cricket sensory data.
	- Use temporal patterns of spikes across a few neurons.
	- Recover the stimulus reconstruction kernel at the coarsest quantization.
	- Demonstrate the presence of additional structure at finer quantizations.
	- Demonstrate non-linear processing in several cells.

